

## Instruction Set Architecture



(a) A seven-bit cell.

Figure 3.1 (continued)

(b) Some possible values in a seven-bit cell.

Figure 3.1
(continued)

(c) Some impossible values in a seven-bit cell.

## Computer Systems fiftu eotrion

## Counting in decimal

| 0 | 7 | 14 | 21 | 28 | 35 |
| ---: | ---: | ---: | ---: | ---: | :---: |
| 1 | 8 | 15 | 22 | 29 | 36 |
| 2 | 9 | 16 | 23 | 30 | 37 |
| 3 | 10 | 17 | 24 | 31 | 38 |
| 4 | 11 | 18 | 25 | 32 | . |
| 5 | 12 | 19 | 26 | 33 | . |
| 6 | 13 | 20 | 27 | 34 | . |

## Computer Systems firtu eortion

## Counting in octal

| 0 | 7 | 16 | 25 | 34 | 43 |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | 10 | 17 | 26 | 35 | 44 |
| 2 | 11 | 20 | 27 | 36 | 45 |
| 3 | 12 | 21 | 30 | 37 | 46 |
| 4 | 13 | 22 | 31 | 40 | . |
| 5 | 14 | 23 | 32 | 41 | . |
| 6 | 15 | 24 | 33 | 42 | . |

## Computer Systems fifth edition

## Counting in base 3

| 0 | 21 | 112 | 210 | 1001 | 1022 |
| ---: | ---: | ---: | ---: | ---: | :---: |
| 1 | 22 | 120 | 211 | 1002 | 1100 |
| 2 | 100 | 121 | 212 | 1010 | 1101 |
| 10 | 101 | 122 | 220 | 1011 | 1102 |
| 11 | 102 | 200 | 221 | 1012 | . |
| 12 | 110 | 201 | 222 | 1020 | . |
| 20 | 111 | 202 | 1000 | 1021 | . |

## Computer Systems fifth edition

## Counting in binary

| 0 | 111 | 1110 | 10101 | 11100 | 100011 |
| ---: | :---: | :---: | :---: | :---: | :---: |
| 1 | 1000 | 1111 | 10110 | 11101 | 100100 |
| 10 | 1001 | 10000 | 10111 | 11110 | 100101 |
| 11 | 1010 | 10001 | 11000 | 11111 | 100110 |
| 100 | 1011 | 10010 | 11001 | 100000 | . |
| 101 | 1100 | 10011 | 11010 | 100001 | . |
| 110 | 1101 | 10100 | 11011 | 100010 | . |


(a) The place values for 10110 (bin).

0 1's place $=0$
1 2's place $=2$
1 4's place $=4$
0 8's place $=0$
1 16's place $=16$
$\overline{22}$ (dec)
(b) Converting 10110 (bin) to decimal.

## Computer Systems

Figure 3.3


## Computer Systems firtu eortion

Figure 3.4

$$
1 \times 2^{4}+0 \times 2^{3}+1 \times 2^{2}+1 \times 2^{1}+0 \times 2^{0}
$$

(a) The binary number 10110 .
$5 \times 10^{4}+8 \times 10^{3}+0 \times 10^{2}+3 \times 10^{1}+6 \times 10^{0}$
(b) The decimal number 58,036 .



## Computer Systems rirfu forion

# Binary addition rules 

$$
\begin{aligned}
& 0+0=0 \\
& 0+1=1 \\
& 1+0=1 \\
& 1+1=10
\end{aligned}
$$



Figure 3.8


- The NEG operation
- Taking the two's complement
- The NOT operation
- Change the I's to 0's and the 0's to I's


## Computer Systems firtu eortion

## The two's complement rule

- The two's complement of a number is I plus its one's complement
- NEG $x=1+$ NOT $x$


## Decimal

$$
\begin{array}{ll}
-7 & 1001 \\
-6 & 1010 \\
-5 & 1011 \\
-4 & 1100 \\
-3 & 1101 \\
-2 & 1110 \\
-1 & 1111
\end{array}
$$

## Computer Systems

| Decimal | Binary |  |  |
| :---: | :---: | :---: | :---: |
| -8 | 1000 | 0 | 0000 |
| -7 | 1001 | 1 | 0001 |
| -6 | 1010 | 3 | 0010 |
| -5 | 1011 |  | 0011 |
| -4 | 1100 | 4 | 0100 |
| -3 | 1101 | 5 | 0101 |
| -2 | 1110 | 6 | 0110 |
| -1 | 1111 | 7 | 0111 |

## Computer Systems

Figure 3.II


## Computer Systems

Figure 3.12


## Computer Systems

Figure 3.13

(a) Breaking the number line in the middle.

(b) Shifting the right part to the left side.

## Computer Systems



## Computer Systems

```
#include <stdio.h>
#include <limits.h>
int main() {
    int n = INT_MAX - 2;
    for (int i = 0; i < 6; i++) {
        printf("n == %d\n", n);
        n++;
        }
        return 0;
}
```

Output
$\mathrm{n}=\mathrm{n}^{2147483645}$
$\mathrm{n}==2147483646$
$\mathrm{n}=2147483647$
$\mathrm{n}==-2147483648$
$\mathrm{n}==-2147483647$
$\mathrm{n}=\mathbf{- 2 1 4 7 4 8 3 6 4 6}$

## Computer Systems firtu eortion

## The status bits

- $N=I$ if the result is negative
$\mathrm{N}=0$ otherwise
- $Z=I$ if the result is all zeros
$Z=0$ otherwise
- $\mathrm{V}=\mathrm{I}$ if a signed integer overflow occurred $\mathrm{V}=0$ otherwise
- $C=I$ if an unsigned integer overflow occurred
$C=0$ otherwise


## Computer Systems

## Addition with a 6-bit cell

Adding two positives:

|  | 000011 |  |  |  |
| :--- | :--- | :--- | :--- | :--- |
|  |  |  | 010110 |  |
| ADD | 010101 |  |  |  |
| $\mathrm{~V}=0$ | 011000 |  |  |  |
| $\mathrm{C}=0$ |  |  | ADD <br> $\mathrm{V}=1$ <br> $\mathrm{~V}=1$ | 100010 |
| $\mathrm{C}=0$ |  |  |  |  |

Adding a positive and a negative:

|  | 000101 |  |  |  |
| :--- | :--- | :--- | :--- | :--- |
|  |  |  | 001000 |  |
| ADD | 110111 |  |  |  |
| $\mathrm{~V}=0$ | 111100 |  | ADD | 111010 |
| $\mathrm{C}=0$ |  |  | $\mathrm{V}=0$ <br> $\mathrm{C}=1$ | 000010 |

Adding two negatives:

|  | 111010 |  |  |
| :--- | :--- | :--- | :--- |
|  |  |  | 100110 |
| ADD | 110111 |  |  |
| $\mathrm{~V}=0$ | 110001 |  |  |
| $\mathrm{C}=1$ |  |  | ADD <br> C <br> $\mathrm{V}=100010$ <br> $\mathrm{~V}=1$ |

## Computer Systems

Figure 3.16

| $\mathbf{p}$ | $\mathbf{q}$ | $\mathbf{p A N D} \mathbf{q}$ | $\mathbf{p}$ | $\mathbf{q}$ | $\mathbf{p}$ OR q | $\mathbf{p}$ | $\mathbf{q}$ | $\mathbf{p X O R} \mathbf{q}$ |
| :--- | :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 0 | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 1 |
| 1 | 0 | 0 | 1 | 0 | 1 | 1 | 0 | 1 |
| 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 0 |

## Computer Systems

| $\mathbf{p}$ | $\mathbf{q}$ | $\mathbf{p A N D} \mathbf{q}$ | $\mathbf{p}$ | $\mathbf{q}$ | $\mathbf{p O R q}$ | $\mathbf{p}$ | $\mathbf{q}$ | $\mathbf{p}$ XOR $\mathbf{q}$ |
| :--- | :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| true | true | true | true | true | true | true | true | false |
| true | false | false | true | false | true | true | false | true |
| false | true | false | false | true | true | false | true | true |
| false | false | false | false | false | false | false | false | false | | (a) HOL6 table for AND. | (b) HOL6 table for OR. |
| :--- | :--- |

## Computer Systems

OperationAND
OR
XOR
NOT
Implies
Transfer
Bit indexInformal descriptionSequential separator
RTL Symbol
$\wedge$ ..... v
$\oplus$
$\neg$$\Rightarrow$$\leftarrow$〈〉\{ \};

## Computer Systems firtu eortion

## RTL specification of $O R$ operation

$c \leftarrow a \vee b ; \mathrm{N} \leftarrow c<0, \mathrm{Z} \leftarrow c=0$

## Computer Systems firtu eortion

Figure 3.19

## Arithmetic shift left (ASL)



$$
\begin{aligned}
& \mathrm{C} \leftarrow r\langle 0\rangle, r\langle 0 . .4\rangle \leftarrow r\langle 1 . .5\rangle, r\langle 5\rangle \leftarrow 0 \\
& \mathrm{~N} \leftarrow r<0, \mathrm{Z} \leftarrow r=0, \mathrm{~V} \leftarrow\{\text { overflow }\}
\end{aligned}
$$

## Arithmetic shift right (ASR)



RTL specification is a problem for the student

Figure 3.21
Rotate left (ROL)


Figure 3.21
(continued)

## Rotate right (ROR)



## Computer Systems firtu eortion

## Counting in hexadecimal

| 0 | 7 | E | I5 | IC | 23 |
| ---: | ---: | ---: | ---: | ---: | :---: |
| I | 8 | F | 16 | ID | 24 |
| 2 | 9 | 10 | 17 | IE | 25 |
| 3 | A | 11 | 18 | IF | 26 |
| 4 | B | 12 | 19 | 20 | . |
| 5 | C | 13 | IA | 21 | . |
| 6 | D | 14 | IB | 22 | . |

## Computer Systems

Figure 3.22

(a) The place values for 8 BE 7 .

$$
\begin{array}{rrr}
7 \times & 1 & =7 \\
14 \times 16 & =224 \\
11 \times 256 & = & 2,816 \\
8 \times 4096 & =\frac{32,768}{35,815}
\end{array}
$$

(b) Converting 8BE7 to decimal.

Figure 3.23

|  | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{2}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{5}$ | $\mathbf{6}$ | $\mathbf{7}$ | $\mathbf{8}$ | $\mathbf{9}$ | $\mathbf{A}$ | $\mathbf{B}$ | $\mathbf{C}$ | $\mathbf{D}$ | $\mathbf{E}$ | $\mathbf{F}$ |
| :--- | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| $\mathbf{0}$ | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 |
| $\mathbf{1}_{-}$ | 16 | 17 | 18 | 19 | 20 | 21 | 22 | 23 | 24 | 25 | 26 | 27 | 28 | 29 | 30 | 31 |
| $\mathbf{2}_{-}$ | 32 | 33 | 34 | 35 | 36 | 37 | 38 | 39 | 40 | 41 | 42 | 43 | 44 | 45 | 46 | 47 |
| $\mathbf{3}_{-}$ | 48 | 49 | 50 | 51 | 52 | 53 | 54 | 55 | 56 | 57 | 58 | 59 | 60 | 61 | 62 | 63 |
| $\mathbf{4}_{-}$ | 64 | 65 | 66 | 67 | 68 | 69 | 70 | 71 | 72 | 73 | 74 | 75 | 76 | 77 | 78 | 79 |
| $\mathbf{5}_{-}$ | 80 | 81 | 82 | 83 | 84 | 85 | 86 | 87 | 88 | 89 | 90 | 91 | 92 | 93 | 94 | 95 |
| $\mathbf{6}_{-}$ | 96 | 97 | 98 | 99 | 100 | 101 | 102 | 103 | 104 | 105 | 106 | 107 | 108 | 109 | 110 | 111 |
| $\mathbf{7 -}_{-}$ | 112 | 113 | 114 | 115 | 116 | 117 | 118 | 119 | 120 | 121 | 122 | 123 | 124 | 125 | 126 | 127 |
| $\mathbf{8}_{-}$ | 128 | 129 | 130 | 131 | 132 | 133 | 134 | 135 | 136 | 137 | 138 | 139 | 140 | 141 | 142 | 143 |
| $\mathbf{Q -}_{-}$ | 144 | 145 | 146 | 147 | 148 | 149 | 150 | 151 | 152 | 153 | 154 | 155 | 156 | 157 | 158 | 159 |
| $\mathbf{A}_{-}$ | 160 | 161 | 162 | 163 | 164 | 165 | 166 | 167 | 168 | 169 | 170 | 171 | 172 | 173 | 174 | 175 |
| $\mathbf{B -}_{-}$ | 176 | 177 | 178 | 179 | 180 | 181 | 182 | 183 | 184 | 185 | 186 | 187 | 188 | 189 | 190 | 191 |
| $\mathbf{C}_{-}$ | 192 | 193 | 194 | 195 | 196 | 197 | 198 | 199 | 200 | 201 | 202 | 203 | 204 | 205 | 206 | 207 |
| $\mathbf{D}_{-}$ | 208 | 209 | 210 | 211 | 212 | 213 | 214 | 215 | 216 | 217 | 218 | 219 | 220 | 221 | 222 | 223 |
| $\mathbf{E}_{-}$ | 224 | 225 | 226 | 227 | 228 | 229 | 230 | 231 | 232 | 233 | 234 | 235 | 236 | 237 | 238 | 239 |
| $\mathbf{F}_{-}$ | 240 | 241 | 242 | 243 | 244 | 245 | 246 | 247 | 248 | 249 | 250 | 251 | 252 | 253 | 254 | 255 |

## Hexadecimal Binary

| 0 | 0000 | 8 | 1000 |
| :--- | :--- | :--- | :--- |
| 1 | 0001 | 9 | 1001 |
| 2 | 0010 | A | 1010 |
| 3 | 0011 | B | 1011 |


| 4 | 0100 | C | 1100 |
| :--- | :--- | :--- | :--- |
| 5 | 0101 | D | 1101 |
| 6 | 0110 | E | 1110 |
| 7 | 0111 | F | 1111 |


| Char | Bin | Hex | Char | Bin | Hex | Char | Bin | Hex | Char | Bin | Hex |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| NUL | 0000000 | 00 | SP | 0100000 | 20 | @ | 1000000 | 40 | , | 1100000 | 60 |
| SOH | 0000001 | 01 | ! | 0100001 | 21 | A | 1000001 | 41 | a | 1100001 | 61 |
| STX | 0000010 | 02 | " | 0100010 | 22 | B | 1000010 | 42 | b | 1100010 | 62 |
| ETX | 0000011 | 03 | \# | 0100011 | 23 | C | 1000011 | 43 | c | 1100011 | 63 |
| EOT | 0000100 | 04 | \$ | 0100100 | 24 | D | 1000100 | 44 | d | 1100100 | 64 |
| ENQ | 0000101 | 05 | \% | 0100101 | 25 | E | 1000101 | 45 | e | 1100101 | 65 |
| ACK | 0000110 | 06 | \& | 0100110 | 26 | F | 1000110 | 46 | f | 1100110 | 66 |
| BEL | 0000111 | 07 | ' | 0100111 | 27 | G | 1000111 | 47 | g | 1100111 | 67 |
| BS | 0001000 | 08 | ( | 0101000 | 28 | H | 1001000 | 48 | h | 1101000 | 68 |
| HT | 0001001 | 09 | ) | 0101001 | 29 | I | 1001001 | 49 | i | 1101001 | 69 |
| LF | 0001010 | 0A | * | 0101010 | 2A | J | 1001010 | 4A | j | 1101010 | 6A |
| VT | 0001011 | OB | + | 0101011 | 2B | K | 1001011 | 4B | k | 1101011 | 6B |
| FF | 0001100 | OC | , | 0101100 | 2 C | L | 1001100 | 4 C | 1 | 1101100 | 6 C |
| CR | 0001101 | OD | - | 0101101 | 2D | M | 1001101 | 4D | m | 1101101 | 6D |
| SO | 0001110 | OE | - | 0101110 | 2E | N | 1001110 | 4E | n | 1101110 | 6E |
| SI | 0001111 | OF | / | 0101111 | 2F | 0 | 1001111 | 4F | $\bigcirc$ | 1101111 | 6F |
| DLE | 0010000 | 10 | 0 | 0110000 | 30 | P | 1010000 | 50 | p | 1110000 | 70 |
| DC1 | 0010001 | 11 | 1 | 0110001 | 31 | Q | 1010001 | 51 | q | 1110001 | 71 |
| DC2 | 0010010 | 12 | 2 | 0110010 | 32 | R | 1010010 | 52 | r | 1110010 | 72 |
| DC3 | 0010011 | 13 | 3 | 0110011 | 33 | S | 1010011 | 53 | s | 1110011 | 73 |
| DC4 | 0010100 | 14 | 4 | 0110100 | 34 | T | 1010100 | 54 | t | 1110100 | 74 |
| NAK | 0010101 | 15 | 5 | 0110101 | 35 | U | 1010101 | 55 | u | 1110101 | 75 |
| SYN | 0010110 | 16 | 6 | 0110110 | 36 | V | 1010110 | 56 | v | 1110110 | 76 |
| ETB | 0010111 | 17 | 7 | 0110111 | 37 | W | 1010111 | 57 | w | 1110111 | 77 |
| CAN | 0011000 | 18 | 8 | 0111000 | 38 | X | 1011000 | 58 | x | 1111000 | 78 |
| EM | 0011001 | 19 | 9 | 0111001 | 39 | Y | 1011001 | 59 | y | 1111001 | 79 |
| SUB | 0011010 | 1A | : | 0111010 | 3 A | Z | 1011010 | 5A | z | 1111010 | 7A |
| ESC | 0011011 | 1B | ; | 0111011 | 3B | [ | 1011011 | 5B | 1 | 1111011 | 7B |
| FS | 0011100 | 1C | < | 0111100 | 3 C | $\backslash$ | 1011100 | 5C |  | 1111100 | 7 C |
| GS | 0011101 | 1D | = | 0111101 | 3D | ] | 1011101 | 5D | \} | 1111101 | 7D |
| RS | 0011110 | 1E | > | 0111110 | 3E |  | 1011110 | 5E | $\sim$ | 1111110 | 7E |
| US | 0011111 | 1F | ? | 0111111 | 3F |  | 1011111 | 5F | DEL | 1111111 | 7F |

Figure 3.25 (continued)

Abbreviations for Control Characters

| NUL | null, or all zeros |
| :--- | :--- |
| SOH | start of heading |
| STX | start of text |
| ETX | end of text |
| EOT | end of transmission |
| ENO | enquiry |
| ACK | acknowledge |
| BEL | bell |
| BS | backspace |
| HT | horizontal tabulation |
| LF | line feed |
| VT | vertical tabulation |


| FF | form feed |
| :--- | :--- |
| CR | carriage return |
| SO | shift out |
| SI | shift in |
| DLE | data link escape |
| DC1 | device control 1 |
| DC2 | device control 2 |
| DC3 | device control 3 |
| DC4 | device control 4 |
| NAK | negative acknowledge |
| SYN | synchronous idle |
| ETB | end of transmission block |


| CAN | cancel |
| :--- | :--- |
| EM | end of medium |
| SUB | substitute |
| ESC | escape |
| FS | file separator |
| GS | group separator |
| RS | record separator |
| US | unit separator |
| SP | space |
| DEL | delete |

## Computer Systems

Figure 3.26

| Unicode Script | Code Point | Glyphs |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
| Arabic | U＋063＿ | ذ | J | j | س | ش | － | ض | b |
| Armenian | U＋054＿ | 2 | 2 | 亿． | $\chi$ | U | 3 | b | c |
| Braille Patterns | U＋287＿ | ： | $\because$ | $\because$ | $\because$ | ： | $\vdots$ | ： | ！： |
| CJK Unified | U＋4EB＿ | 京 | 但 | 亲 | 毫 | 毫 | 熟 | 亶 | 廉 |
| Cyrillic | U＋041＿ | A | Б | B | $\Gamma$ | Д | E | Ж | 3 |
| Egyptian Hieroglyphs | U＋1300＿ | 易 | 賏 | 通 | 曷 |  | 䫃 | 䱊 | 暘 |
| Emoticons | U＋1F61＿ | ； | $\Theta$ | $\bigcirc$ | （3） | O | （2） | （） | － |
| Hebrew | U＋05D＿ | $\aleph$ | $\geq$ | $\lambda$ | 7 | $\cdots$ | 1 | i | $\Pi$ |
| Basic Latin（ASCII） | U＋004＿ | ＠ | A | B | C | D | E | F | G |
| Latin－1 Supplement | U＋00E＿ | à | á | â | a | ä | a | æ | ç |

## Computer Systems

## UTF-8 encoding

| Bits | First Code Point | Last Code Point | Byte 1 | Byte 2 | Byte 3 |
| :--- | :--- | :--- | :--- | :--- | :--- | Byte 4

# Floating point representation 

## Computer Systems


(a) The place values for 101.011 (bin).
$11 / 8$ 's place $=0.125$
1 1/4's place $=0.25$
$0 \frac{1}{2}$ 's place $=0.0$
1 1's place $=1.0$
0 2's place $=0.0$
1 4's place $=4.0$ 5.375 (dec)
(b) Converting 101.011 (bin) to decimal.

## Computer Systems

Figure 3.29

$$
1 \times 2^{2}+0 \times 2^{1}+1 \times 2^{0}+0 \times 2^{-1}+1 \times 2^{-2}+1 \times 2^{-3}
$$

(a) The binary number 101.011.
$5 \times 10^{2}+0 \times 10^{1}+6 \times 10^{0}+7 \times 10^{-1}+2 \times 10^{-2}+1 \times 10^{-3}$
(b) The decimal number 506.721.

## Computer Systems

6.5859375
$6(\mathrm{dec})=110(\mathrm{bin})$
(a) Convert the whole part

|  | .5859375 |
| :--- | :--- |
| 1 | .171875 |
| 0 | .34375 |
| 0 | .6875 |
| 1 | .375 |
| 0 | .75 |
| 1 | .5 |
| 1 | .0 |

(b) Convert the fractional part

## Computer Systems rifтu eortion

## Normalized

- Leading I on the left of the binary point
- $6.5859375(\mathrm{dec})=110.1001011$ (bin)
- Normalized scientific notation:
I.101001011 $\times 2^{2}$




## Computer Systems fifth edition

## The hidden bit

- Normalized scientific notation always has I to the left of the binary point
- So, do not store it
- Increases precision in the significand
- Floating point unit inserts hidden bit before doing computation
- Floating point unit removes leading I from significand before storing result


## Computer Systems

Figure 3.33

## Decimal Excess 3 Two's Complement

| -4 |  |
| ---: | :---: |
| -3 | 000 |
| -2 | 001 |
| -1 | 010 |
| 0 | 011 |
| 1 | 100 |
| 2 | 101 |
| 3 | 110 |
| 4 | 111 |

## Computer Systems firtu eortion

Figure 3.34

## Round to nearest

## Ties to even

| Decimal | Decimal <br> Rounded | Binary | Binary <br> Rounded |
| :--- | :--- | :--- | :--- |
| 23.499 | 23 | 1011.011 | 1011 |
| 23.5 | 24 | 1011.1 | 1100 |
| 23.501 | 24 | 1011.101 | 1100 |
| 24.499 | 24 | 1100.011 | 1100 |
| 24.5 | 24 | 1100.1 | 1100 |
| 24.501 | 25 | 1100.101 | 1101 |

## Computer Systems firtu eortion

## Special value

- Zero
- Exponent field all 0's
- Significand all 0's
- There is a +0 and a -0


## Computer Systems



## Special Value

Zero
Denormalized
Infinity
Not a number

## Exponent

All zeros
All zeros
All ones
All ones

Significand
All zeros
Nonzero
All zeros
Nonzero

## Computer Systems firtu eortion

## Special value

- Infinity
- Exponent field all I's
- Significand all 0's
- There is a $+\infty$ and $\mathrm{a}-\infty$
- Produced by operation that gives result in overflow region


## Computer Systems firtu eortion

## Special value

- Not a Number ( $\mathrm{NaN)}$
- Exponent field all I's
- Significand nonzero
- Produced by illegal math operations


## Computer Systems fieru sortion

## Special value

- Denormalized number
- Exponent field all 0's
- Significand nonzero
- Hidden bit is assumed to be 0 instead of I
- If the exponent is stored in excess $n$ for normalized numbers, it is stored in excess $n-1$ for denormalized numbers


# Computer Systems 

Figure 3.37

## Normalized



Denormalized


## Computer Systems

|  | Binary | Scientific Notation | Decimal |
| :--- | :--- | :--- | :--- |
| Not a number | 1111 nonzero |  |  |
| Negative infinity | 11110000 |  | $-\infty$ |
| Negative | 11101111 | $-1.1111 \times 2^{3}$ | -15.5 |
| normalized | 11101110 | $-1.1110 \times 2^{3}$ | -15.0 |
|  | $\ldots$ | $\ldots$ | $\ldots$ |
|  | 10110001 | $-1.0001 \times 2^{0}$ | -1.0625 |
|  | 10110000 | $-1.0000 \times 2^{0}$ | -1.0 |
|  | 10101111 | $-1.1111 \times 2^{-1}$ | -0.96875 |
|  | $\ldots$ | $\ldots$ | $\ldots$ |
|  | 10010001 | $-1.0001 \times 2^{-2}$ | -0.265625 |
|  | 10010000 | $-1.0000 \times 2^{-2}$ | -0.25 |
|  | 10001111 | $-0.1111 \times 2^{-2}$ | -0.234375 |
| Negative | 10001110 | $-0.1110 \times 2^{-2}$ | -0.21875 |
| denormalized | $\ldots$ | $\ldots$ | $\ldots$ |
|  | 10000010 | $-0.0010 \times 2^{-2}$ | -0.03125 |
|  | 10000001 | $-0.0001 \times 2^{-2}$ | -0.015625 |
|  | 10000000 |  | -0.0 |

## Computer Systems

| Positive zero | 00000000 |  | +0.0 |
| :--- | :--- | :--- | :--- |
| Positive | 00000001 | $0.0001 \times 2^{-2}$ | 0.015625 |
| denormalized | 00000010 | $0.0010 \times 2^{-2}$ | 0.03125 |
|  | $\ldots$ | $\ldots$ | $\ldots$ |
|  | 00001110 | $0.1110 \times 2^{-2}$ | 0.21875 |
|  | 00001111 | $0.1111 \times 2^{-2}$ | 0.234375 |
| Positive | 00010000 | $1.0000 \times 2^{-2}$ | 0.25 |
| normalized | 00010001 | $1.0001 \times 2^{-2}$ | 0.265625 |
|  | $\ldots$ | $\ldots$ | $\ldots$ |
|  | 00101111 | $1.1111 \times 2^{-1}$ | 0.96875 |
|  | 00110000 | $1.0000 \times 2^{0}$ | 1.0 |
|  | 00110001 | $1.0001 \times 2^{0}$ | 1.0625 |
|  | $\ldots$ | $\ldots$ | $\ldots$ |
|  | 01101110 | $1.1110 \times 2^{3}$ | 15.0 |
|  | 01101111 | $1.1111 \times 2^{3}$ | 15.5 |
| Positive infinity | 01110000 |  | $+\infty$ |
| Not a number | 0111 nonzero |  |  |
|  |  |  |  |

## Computer Systems

Figure 3.39

## IEEE 754 floating point



## Computer Systems fieru eortion

## Single precision

- C type: float
- Exponent: 8-bit cell
- Excess 127 representation
- Excess 126 for denormalized numbers
- Exponent: 8-bit cell
- Significand: 23-bit cell


## Computer Systems rifтu eortion

## Double precision

- C type: double
- Exponent: Il-bit cell
- Excess 1023 representation
- Excess 1022 for denormalized numbers
- Exponent: 8-bit cell
- Significand: 52-bit cell

